

Memory Log – Algebraic Process Calculi

SoSe 2023

The following might not be 100% accurate, the process calculus described here is for the most part the “Language of Temporal Ordering Specification”

Defintions

Approximately 10 minutes are given to read and understand the following definitions.

L is the set of action labels. Let $a \in L \cup \{i\}$, $g \in L \cup \{i, \delta\}$, $a_i, b_i, g_i \in L (0 \leq i < n)$ be the set of actions.

We define the set of process constants as

$$C ::= P$$

Where P is the set of processes defined by

	$P ::=$	exit
		stop
actions		$a; P$
enabling		$P \gg P$
disabling		$P[> P$
hiding		hide $[a_0, \dots a_n]$ in P
renaming		$P[a_0/b_0 \dots a_n/b_n]$
parallel composition		$P \mid [a_0, \dots a_n] \mid P$
constants		C

The transitions are given by

1. **exit** $\xrightarrow{\delta}$ **stop**
2. $a; P \xrightarrow{a} P$
3. if $C ::= P$, $P \xrightarrow{a} P'$, then $C \xrightarrow{a} P'$
4. if $P_1 \xrightarrow{a} P'_1$, $a \neq \delta$ then $P_1 \gg P_2 \xrightarrow{a} P'_1 \gg P_2$
5. if $P_1 \xrightarrow{\delta} P_2$ then $P_1 \gg P_2 \xrightarrow{i} P_2$
6. if $P_2 \xrightarrow{a} P'_2$ then $P_1[> P_2 \xrightarrow{a} P'_2$
7. if $P_1 \xrightarrow{a} P'_1$, $a \neq \delta$ then $P_1[> P_2 \xrightarrow{a} P'_1[> P_2$
8. if $P_1 \xrightarrow{\delta} P'_1$ then $P_1[> P_2 \xrightarrow{i} P_2$

9. if $P_2 \xrightarrow{a} P'_2$, $a \neq \delta$ then $P_1[> P_2 \xrightarrow{a} P_1[> P'_2]$
10. if $P_2 \xrightarrow{\delta} P'_2$ then $P_1[> P_2 \xrightarrow{\delta} P'_2]$
11. if $P_1 \xrightarrow{a} P'_1$, $a \notin \{a_0, \dots, a_n\}$ then $\text{hide } [a_0, \dots, a_n] \text{ in } P_1 \xrightarrow{a} \text{hide } [a_0, \dots, a_n] \text{ in } P'_1$
12. if $P_1 \xrightarrow{a} P'_1$, $a \in \{a_0, \dots, a_n\}$ then $\text{hide } [a_0, \dots, a_n] \text{ in } P_1 \xrightarrow{i} \text{hide } [a_0, \dots, a_n] \text{ in } P'_1$
13. (renaming rules, they behave as expected and where not needed in the exam)
14. if $P_1 \xrightarrow{a} P'_1$, $a \notin \{a_0, \dots, a_n, \delta\}$ then $P_1 \mid [a_0, \dots, a_n] \mid P_2 \xrightarrow{a} P'_1 \mid [a_0, \dots, a_n] \mid P_2$
15. if $P_1 \xrightarrow{a} P'_1$, $P_2 \xrightarrow{a} P'_2$, $a \in \{a_0, \dots, a_n, \delta\}$ then $P_1 \mid [a_0, \dots, a_n] \mid P_2 \xrightarrow{a} P'_1 \mid [a_0, \dots, a_n] \mid P'_2$

1st Prompt

Describe the transitions of the following process

$$(a; \mathbf{exit}) \gg P$$

where $P := i; P$.

2nd Prompt

Describe the transitions of the following process

$$(a; \mathbf{exit})[> P$$

keeping the definition of P from the first prompt.

3rd Prompt

Change the calculus in such a way that ‘enabling’ a process does not require performing an ‘internal’ (i.e. \xrightarrow{i}) transition.

4th Prompt

Let

$$B := a; \mathbf{exit}$$

$$H := \text{hide } [a] \text{ in } B \mid [a] \mid (B \mid [a] \mid B).$$

Describe the transitions.